Extending Stone Duality to Relations

M. Andrew Moshier¹ Achim Jung² Alexander Kurz³ July 2018 Chapman University University of Birmingham University of Leicester

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Investigate

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- Clearly, this will require us to add something to the algebraic side.
- We know what to do in specific cases: Proximity lattices (Smyth, Jung/Sünderhauf), proximity lattices with "negation" (M).

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First step: Relations

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 Proximity lattices are distributive lattices equipped with particular sorts of relations.



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- The dual structures (compact pospaces) are obtained as certain quotients of the underlying dual Priestley spaces (a Stone space is a Priestley space with discrete order).
- To generalize this, we need to understand how relations generally behave under natural dualities.



Relations Three Ways

Spans: Span

For posets X and Y, a span from X to Y is a pair of monotonic functions

$$X \xleftarrow{p} P \xrightarrow{q} Y$$

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- Horizontal composition is defined by commas (the order analogue of pullback).
- ► A 2-morphism from span $X \xleftarrow{p} R \xrightarrow{q} Y$ to $X \xleftarrow{p'} R' \xrightarrow{q'} Y$ is a monotonic function $f: R \to R'$ making the obvious triangles commute.

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► A 2-morphism from cospan $X \xrightarrow{j} C \xleftarrow{k} Y$ to cospan $X \xrightarrow{j'} C' \xleftarrow{k'} Y$ is a monotonic function $f: C \to C'$ making the obvious triangles commute.

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Relations three ways

Weakening relations: WRel

► For posets X and Y, a weakening relation is a monotonic map $R: X^{\partial} \times Y \to 2$. Equivalently, identifying with the co-kernel $R = \{(x, y) \mid R(x, y) = 1\}$: $\underline{x \leq_X x' \quad x' R y' \quad y' \leq_X y}$

x R y

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 Horizontal composition is defined by the usual relation product.



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$$\frac{x \leq_X x' \quad x' \mathrel{R} y' \quad y' \leq_X y}{x \mathrel{R} y}$$

- Horizontal composition is defined by the usual relation product.
- A 2-morphism between weakening relations is simply comparison point-wise.



How these are related?

Weakening relations, spans and cospans form 2-categories. The 2 cells are related via the following functors.

- $R \in WRel$, determines
 - a span graph(R) by restricting projections
 - a cospan collage(R) by taking the least order on X ⊎ Y containing ≤_X, ≤_Y and R

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Weakening relations, spans and cospans form 2-categories. The 2 cells are related via the following functors.

• $X \xleftarrow{p} R \xrightarrow{q} Y$ determines

- ▶ a weakening relation $\operatorname{rel}_{s}(p,q)$ by $(x, y) \in \operatorname{rel}_{s}(p,q)$ iff $\exists r \in R, x \leq p(r)$ and $q(r) \leq y$
- ▶ a cospan cocomma(*p*, *q*) by taking the cocomma of (*p*, *q*).



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 - ► a cospan cocomma(p, q) by taking the cocomma of (p, q).
- $X \xrightarrow{j} C \xleftarrow{k} Y$ determines
 - a weakening relation $\operatorname{rel}_c(j,k)$ by (x,y) iff $j(x) \le k(y)$
 - a span comma(j, k) by taking the comma of (j, k).

How are these related?

We have three 2-categories: Span, Cospan and WRel. We already described the hom categories: Span(X, Y), Cospan(X, Y) and WRel(X, Y).

- Composition of spans is defined by a comma
- Composition of cospans is defined by a cocomma
- Composition of weakening relations is defined by relational product: *R*; *S*(*x*, *y*) = ∨_{*y*∈*Y*} *R*(*x*, *y*) ∧ *S*(*y*, *z*).



How are these related?

The constructions rel_s , rel_c , graph, collage, comma, cocomma are 2-functors:

- $\operatorname{rel}_{s}(X, Y) \dashv \operatorname{graph}(X, Y);$
- ▶ $\operatorname{rel}_{s}(X, Y) \circ \operatorname{graph}(X, Y) \cong \operatorname{WRel}(X, Y)$
- $\operatorname{rel}_c(X, Y) \dashv \operatorname{collage}(X, Y);$
- ▶ $\operatorname{rel}_c(X, Y) \circ \operatorname{collage}(X, Y) \cong \operatorname{WRel}(X, Y);$
- $\operatorname{cocomma}(X, Y) \dashv \operatorname{comma}(X, Y)$
- comma(X, Y) \cong graph(X, Y) \circ rel_c(X, Y).
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- comma(X, Y) \cong graph(X, Y) \circ rel_c(X, Y).
- $\operatorname{cocomma}(X, Y) \cong \operatorname{collage}(X, Y) \circ \operatorname{rel}_{s}(X, Y).$
- These facts hold analogously in PoSpace, the category of topological spaces with closed partial orders with respect to continuous monotonic functions.

Extending to algebras and topological structures

Suppose \mathcal{A} is a class of ordered algebras (algebras with a partial order in which operations are monotone). Let $\overline{\mathcal{A}}$ denote the category of \mathcal{A} -algebra spans in \mathcal{A} with weakening poset reducts.

For example, DLat is the category of bounded distributive lattices with morphisms that are relations satisfying:

- $x \le x' R y' \le y$ implies x R y
- ► 0 *R* 0
- ▶ 1 *R* 1
- $x_0 R y_0$ and $x_1 y_1$ implies $x_0 \wedge x_1 R y_0 \wedge y_1$
- $x_0 R y_0$ and $x_1 R y_1$ implies $x_0 \lor x_1 R y_{\lor} y_1$.



Main point

Theorem

- ► *DL* is (dually equivalent to Priestley.
- Pos is dually equivalent to Stone(DLat)
- SLat is dually equivalent to Stone(SLat).

Proof idea:

► A span $X \xleftarrow{p} R \xrightarrow{q} Y$ in any of the categories mentioned here dualizes to $2^X \xrightarrow{2^p} 2^R \xleftarrow{2^q} 2^Y$ in Priestley.



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- But this transfer preserves the weakening property in each case.
- The correspondence of spans and cospans allows the cospan in the dual category to be tranfered into a span.



So far, we are still in the realm of Stone spaces.

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- By splitting idempotents (below identity) in the algebraic relational categories, we dualize to obtain suitable (pre)congruences in the corresponding topological categories.

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Happy Birthday Dana. Thanks Klaus.